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# *Modelling and Performance Analysis of WMAC Protocols*

## *Overview*

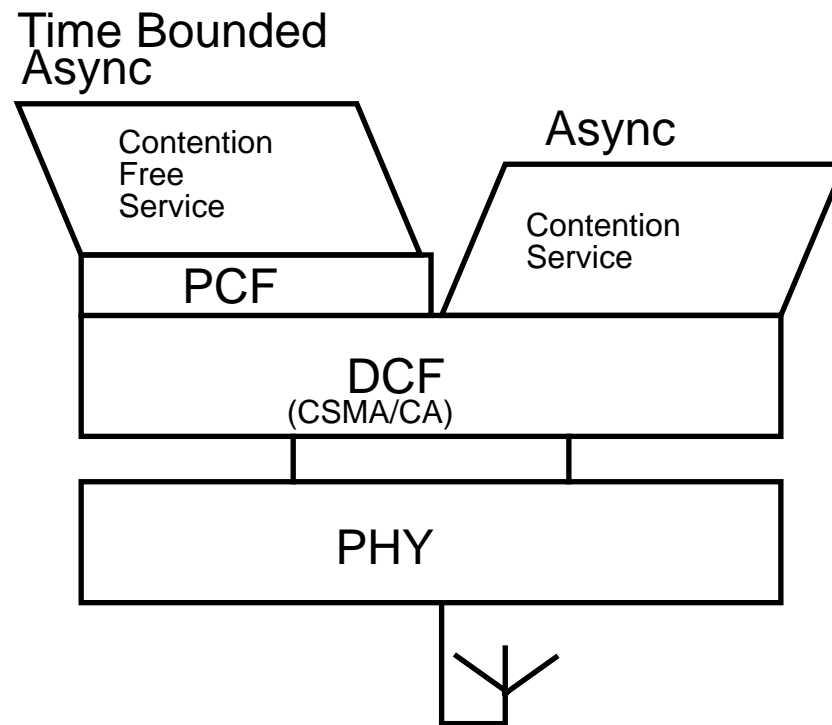
1. DFWMAC
2. Simulation tool
3. DFWMAC model
4. Simulation results
5. Conclusions

## *DFWMAC*

- **Preferred MAC protocol in the IEEE 802.11 standardisation group**  
⇒ *Draft standard*
- **Consists of two parts**
  - ◆ *Basic access method (pseudo CSMA/CA)*
  - ◆ *Contention Free Service support*
- **Two working modes**
  - ◆ *Ad-Hoc*
  - ◆ *Infrastructure (can overlap with Ad-Hoc)*
- **Supported services**
  - ◆ *Asynchronous Data Service*
  - ◆ *Contention Free Service (optional)*
    - Asynchronous Contention Free
    - Time Bounded Service (voice)

## DFWMAC (cont.)

### Service Model



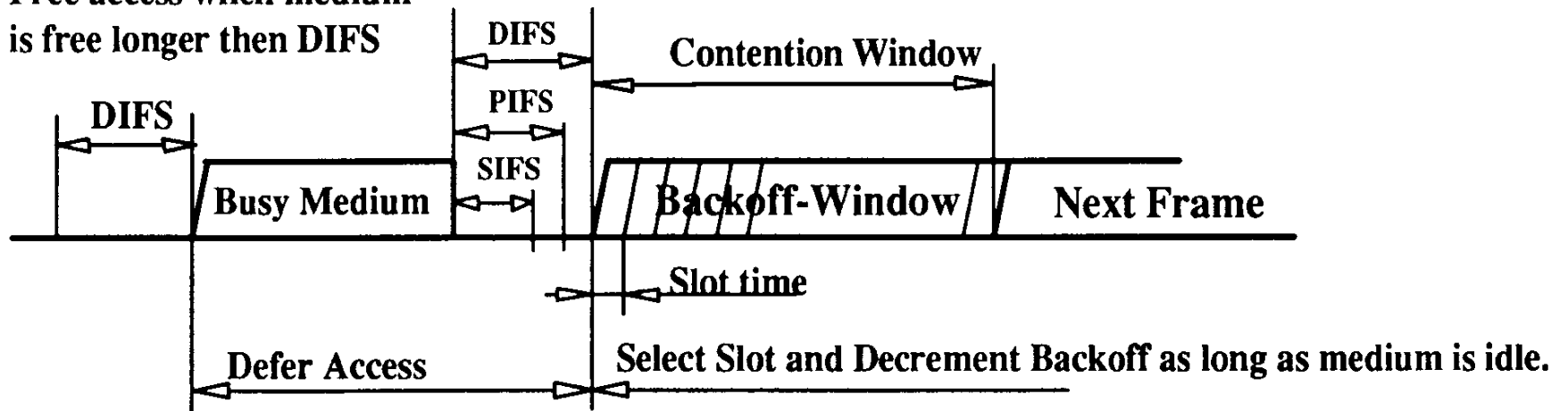
## DFWMAC (cont.)

### Basic access method

- **Distributed Access Protocol**

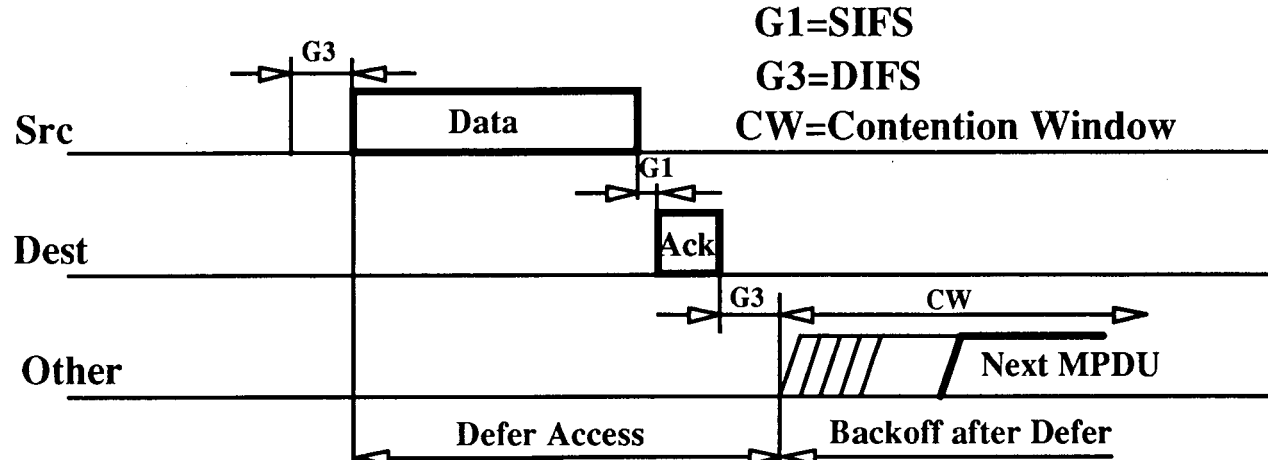
- ♦ MAC level recovery (immediate Ack for unicast transmission)
- ♦ optional RTS/CTS (provide virtual circuit - hidden node problem)
- ♦ different priority levels (IFS - needed for Ack or PCF)

Free access when medium  
is free longer then DIFS

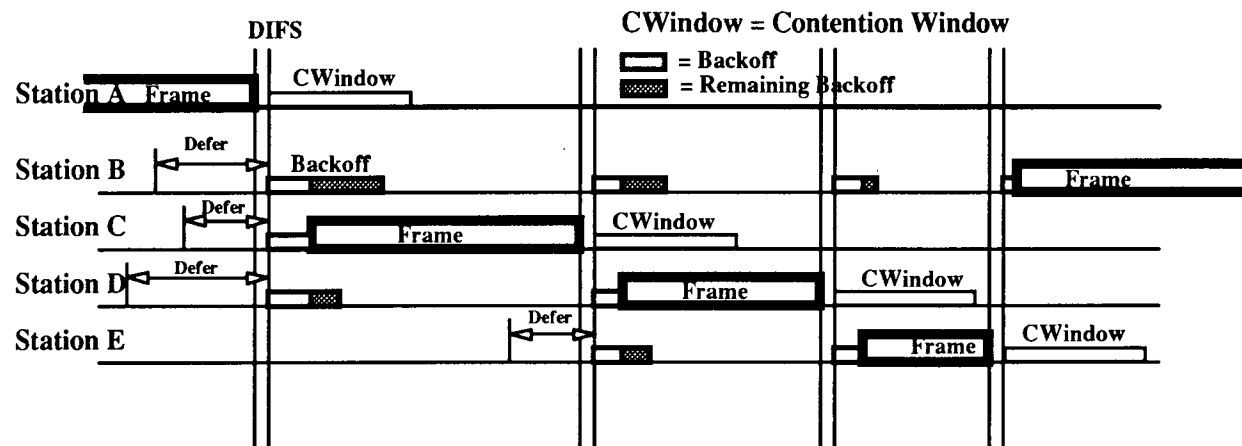


## DFWMAC (cont.)

Access procedure with immediate Ack

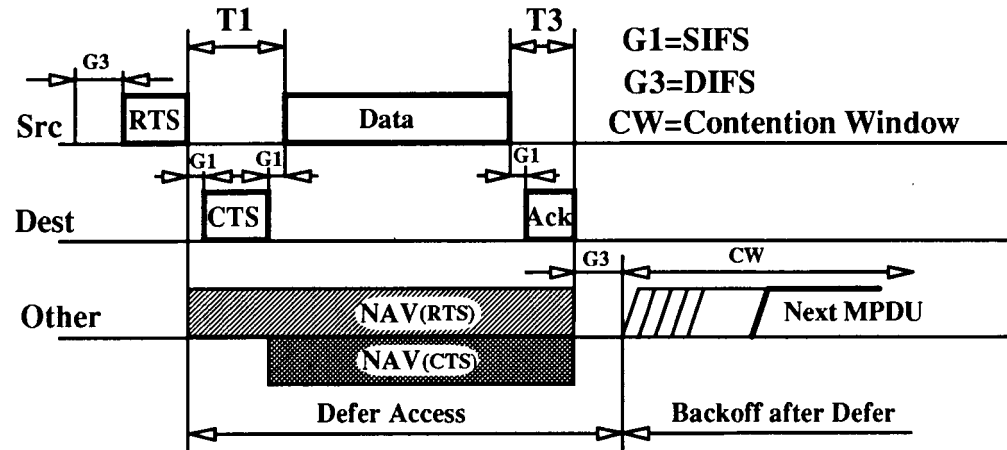


Access Backoff procedure

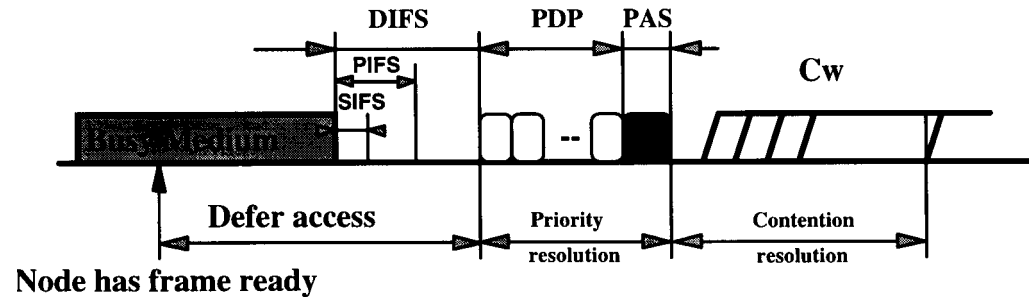


## DFWMAC (cont.)

### RTS/CTS extension

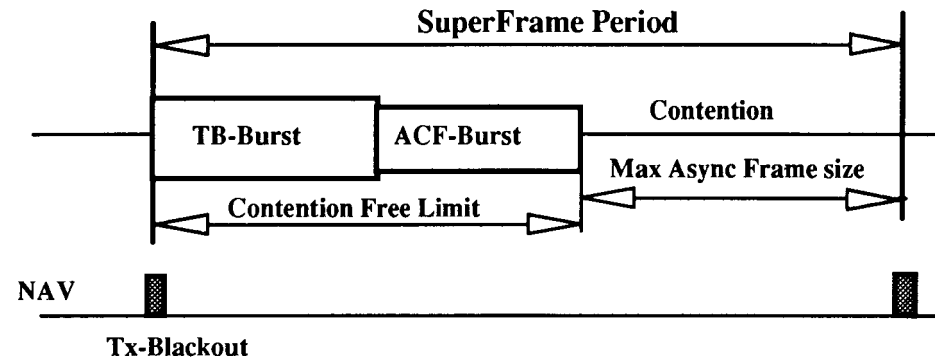


### Active priority signalling (APS)



## DFWMAC (cont.)

### Contention Free Service



- **Registration for time bounded and asynchronous contention free transmission in the contention period**
- **polling of terminals by a control station**
- **use of shorter IFS for high priority traffic**

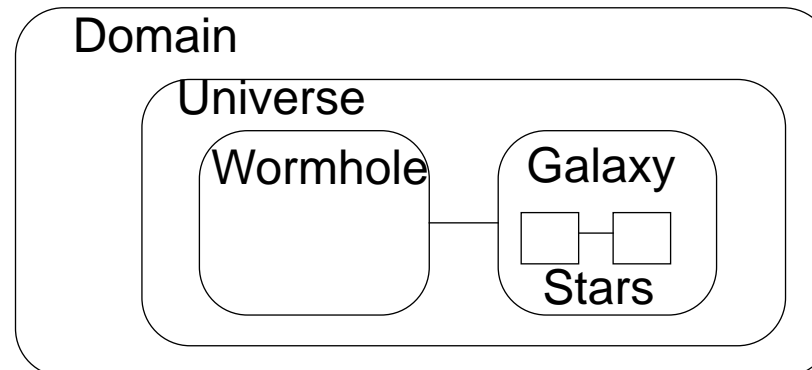
## Simulation tool

- available tools: PTOLEMY, OPnet, SES/workbench, SMURPH (LANSF), CSIM

⇒ PTOLEMY

- general purpose simulation tool
- many connectable simulation environments (discrete event, synchronous data flow, concurrent processes) for several simulation types
- object oriented
- relative simple modelling (rich library)
- graphical interface or Tcl-based interpreter
- free source code
- lots of users (newsgroup, maintenance)

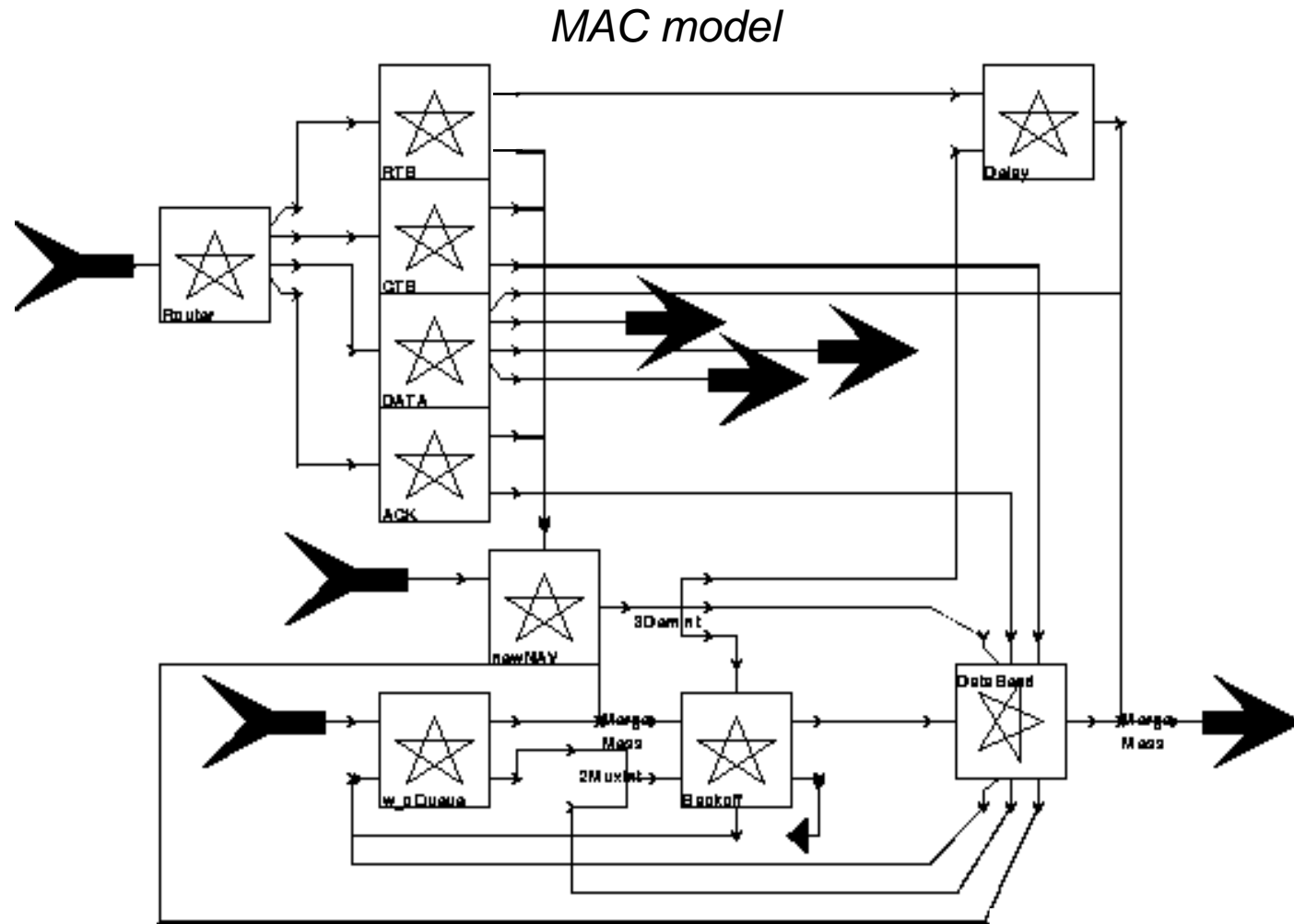
- PTOLEMY model structure





## DFWMAC model

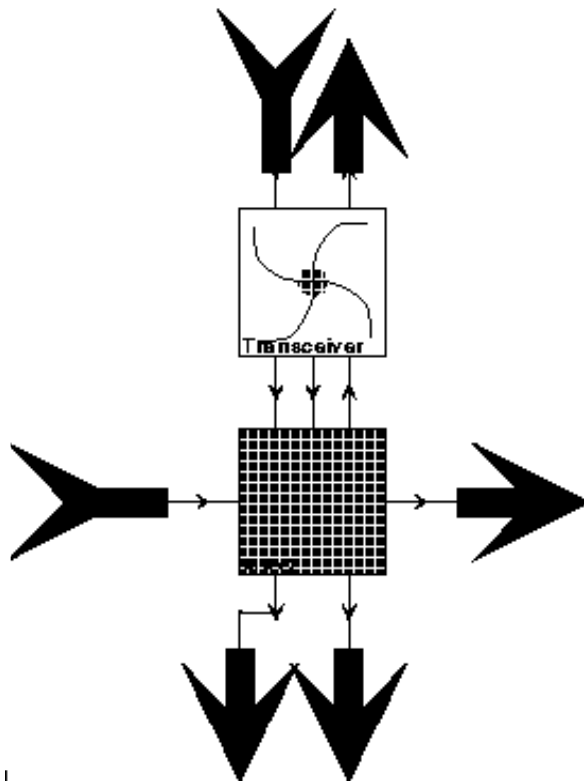
- only DCF (CSMA/CA with Ack and RTS/CTS implemented)



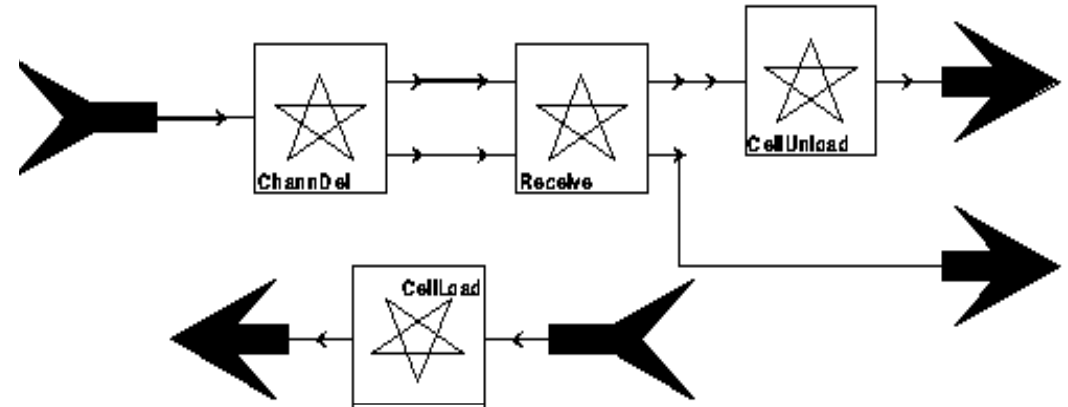
## DFWMAC model (cont.)

### MAC model and transceiver

MAC+Transceiver

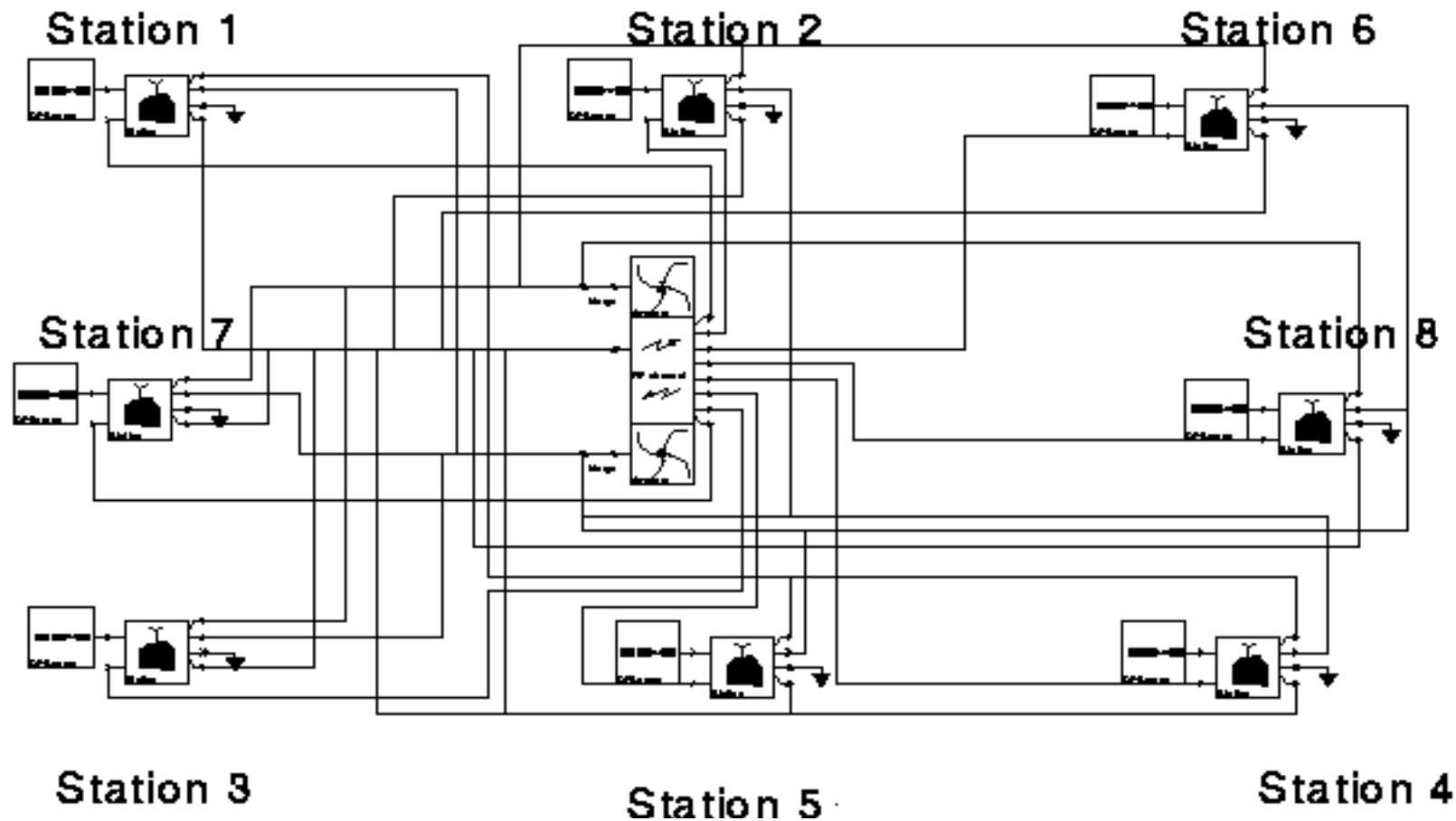


Transceiver



## DFWMAC model (cont.)

Simulation model with 8 stations

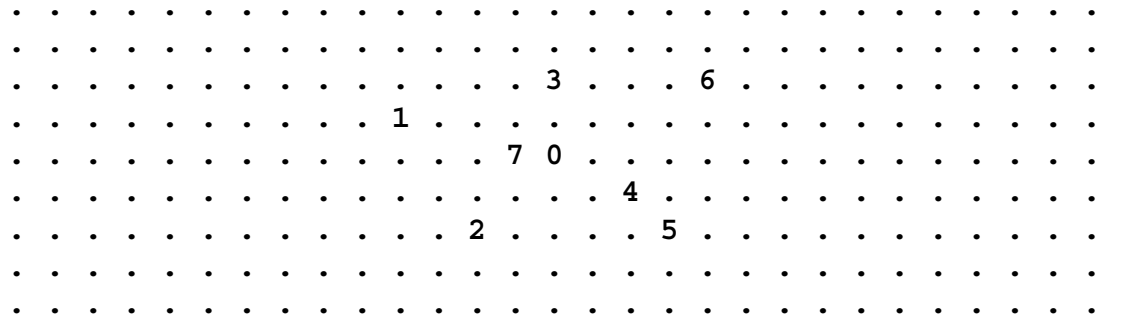


## DFWMAC model (cont.)

### Channel model

- propagation delay

⇒ position of stations (one dot = 10 meter)



- transmission rate (2 Mbit/s)
- attenuation coefficient (2.0)
- Error Probability (Poisson process, mean value  $10^{-4}$ )
- currently no moving stations
- currently no co-channel, adjacent channel interference

## *Simulation results*

- **Simulation of**
  - ◆ *Throughput*
  - ◆ *Access delay*
- **Parameter of simulation**
  - ◆ *Channel Access Mechanism*
    - pure CSMA/CA with/without binary exponential backoff
    - CSMA/CA with Ack (DFWMAC) with/without RTS/CTS
    - CSMA/CA with Ack (DFWMAC) with/without RTS/CTS and APS
    - CSMA/CA with Ack (DFWMAC) with modified backoff algorithm no APS
  - ◆ *Contention Window Size*
  - ◆ *Load pattern:*

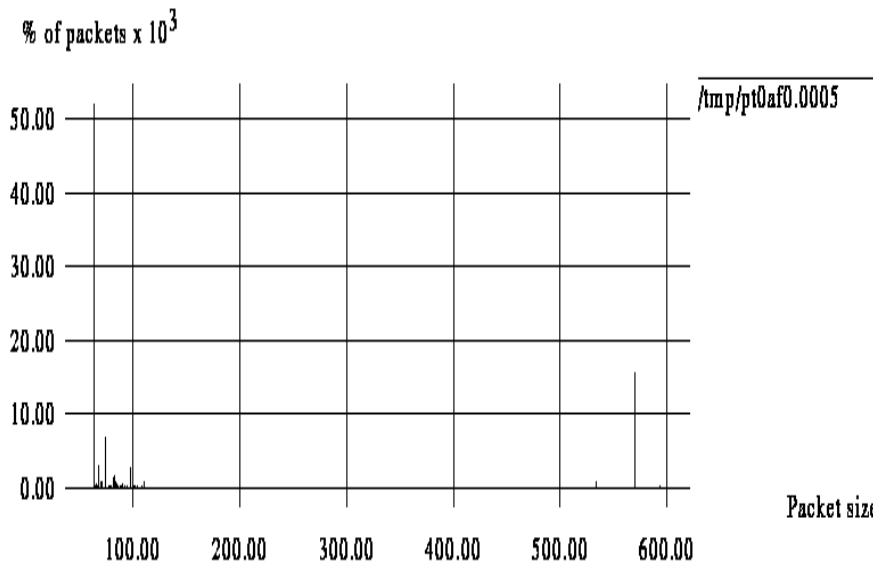
$$\text{Load} = \frac{1}{(\text{InterPacketTime} + \text{TransmissionTime})} \cdot \text{PacketSize}$$
  - ◆ *Packet Size*
    - fixed packet size of 53, 64, 256, 512, 1024 bytes
    - variable packet size (Leland trace)
  - ◆ *no RTS/CTS (RTS/CTS depending on packet size)*

## Simulation results

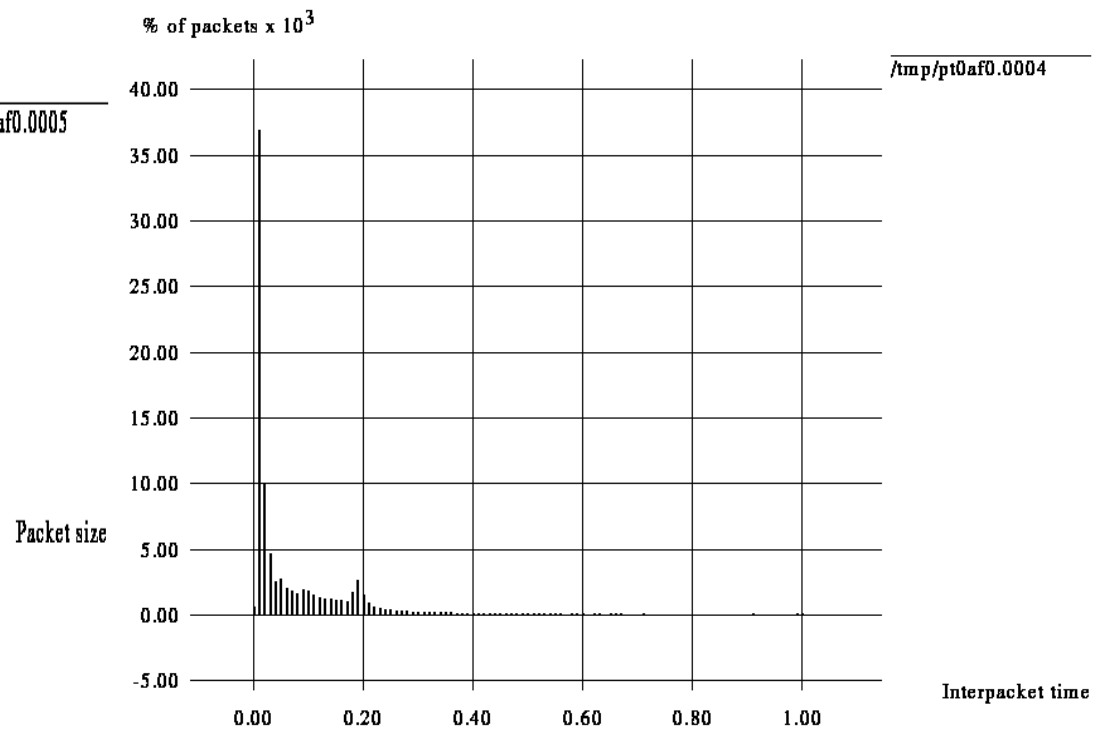
### • Leland trace

- ◆ measurement on a 10BaseT Ethernet at Bellcore Morristown, October 1989
- ◆ 1 million packet arrivals
- ◆ accuracy roughly 10 msec
- ◆ Note: limitation to Ethernet packet sizes (64 to 1518 bytes)

Packet size Leland trace

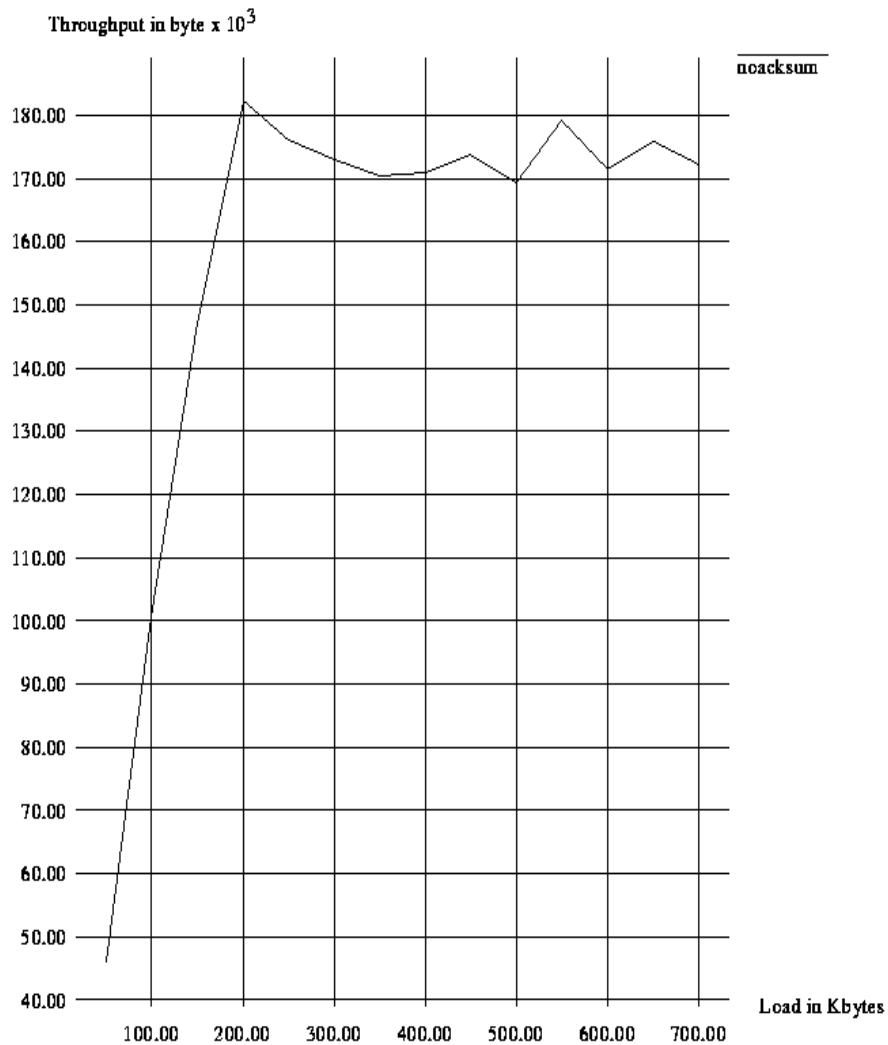


Interpacket Time Leland trace

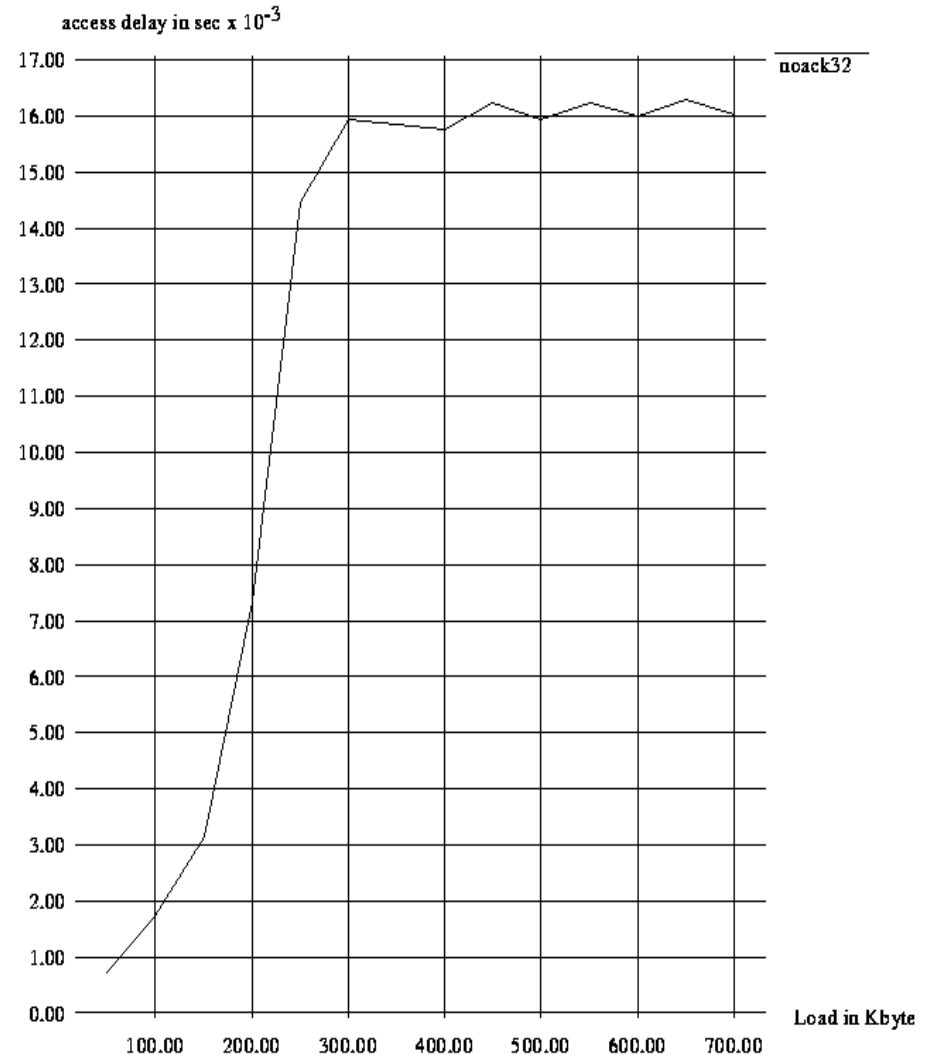


## Simulation results (cont.)

CSMA without ACK, Leland trace, CW=32

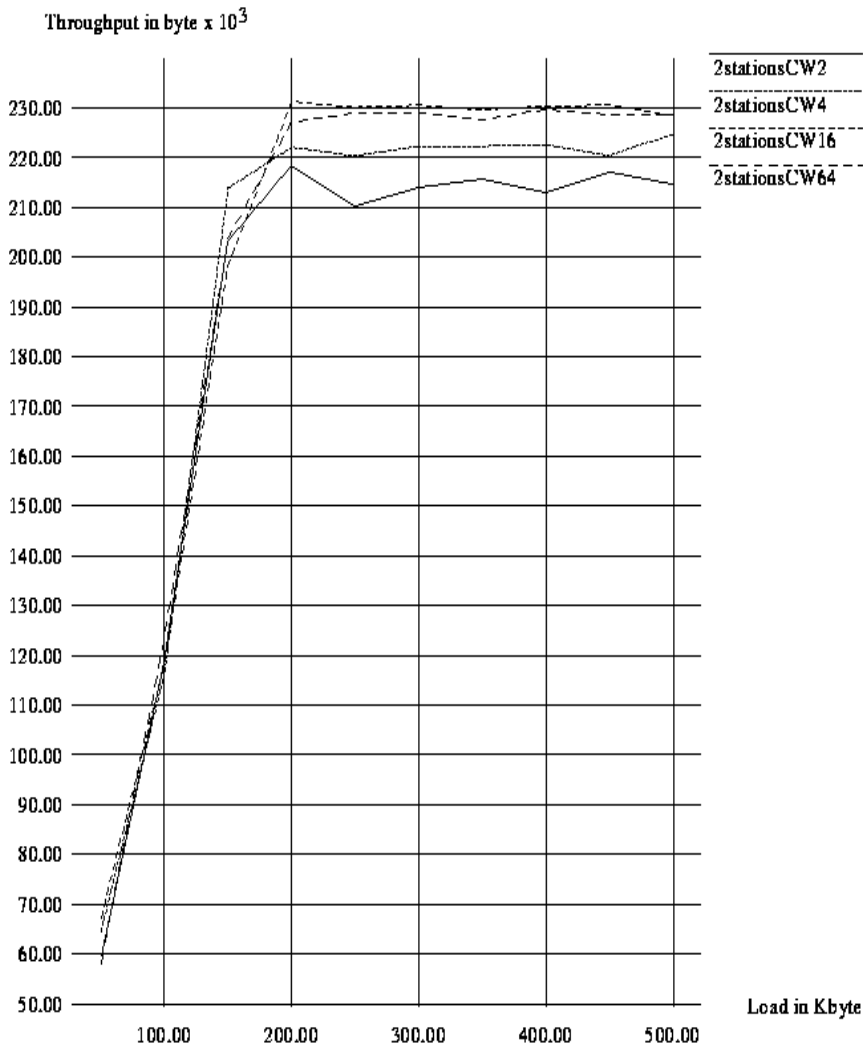


CSMA without ACK, Access delay

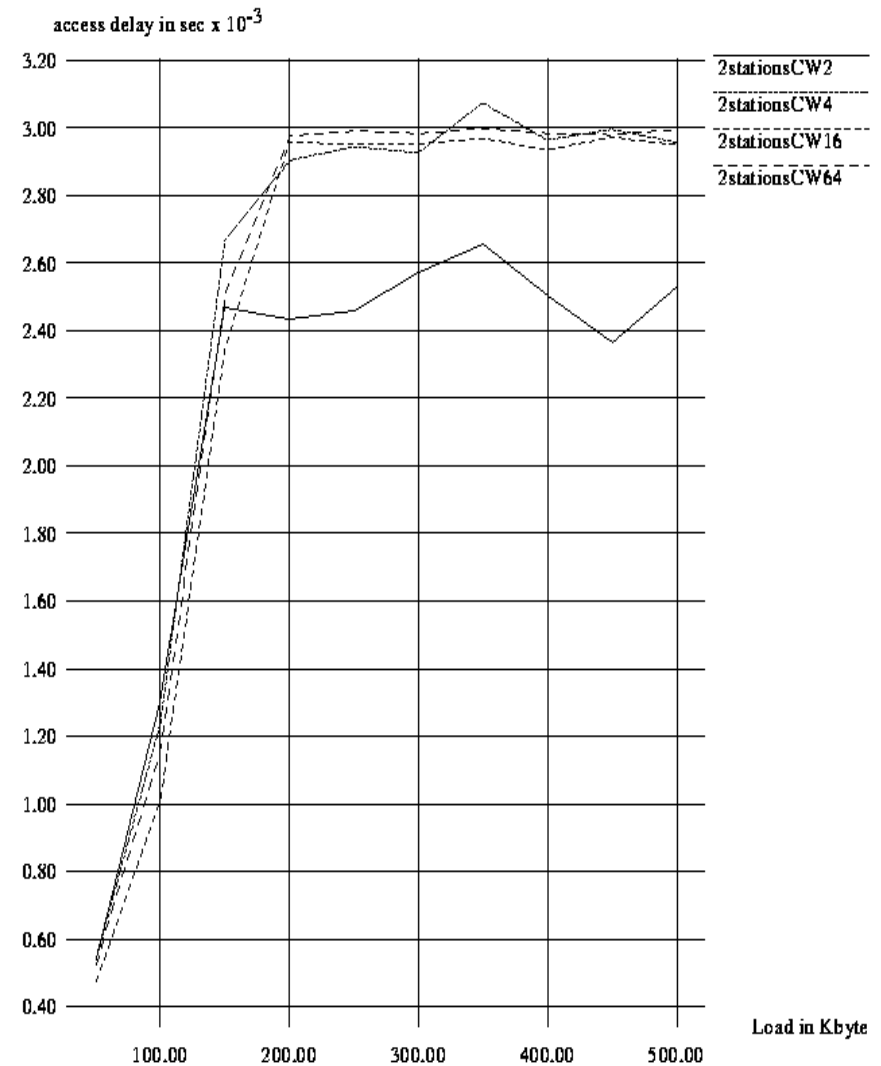


## Simulation results (cont.)

2 stations, Throughput for diff. CW sizes



2 stations, Access delay for diff. CW sizes

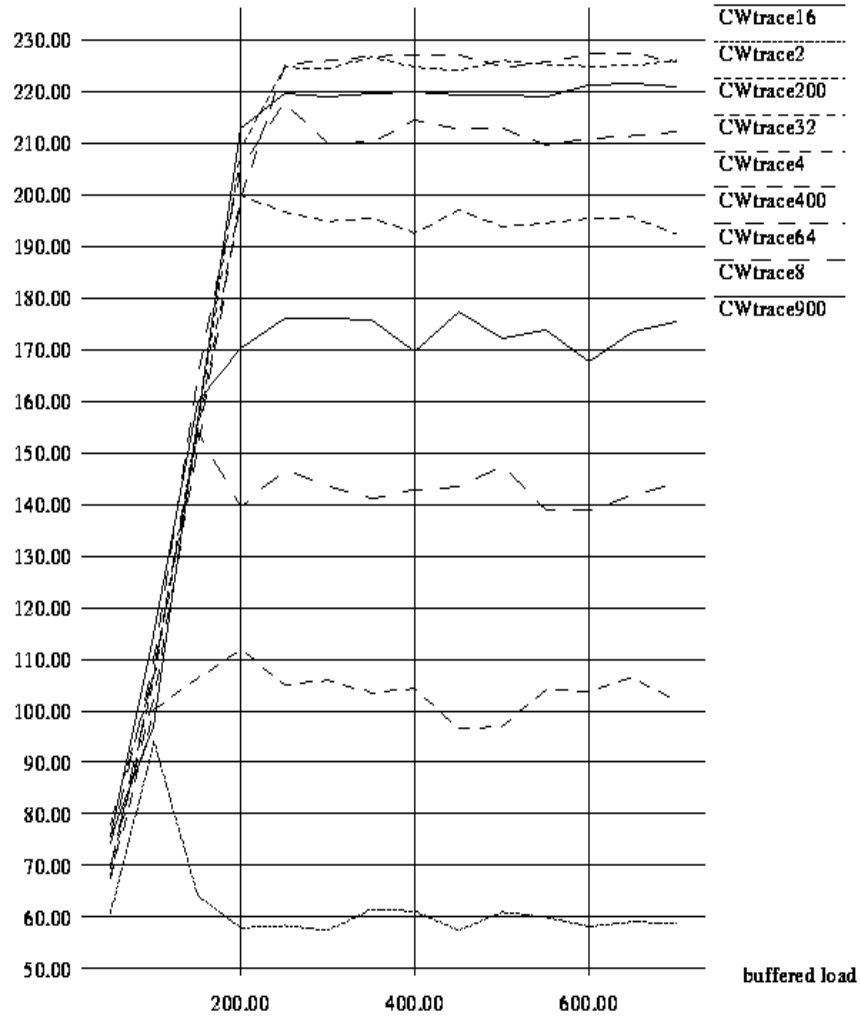




## Simulation results (cont.)

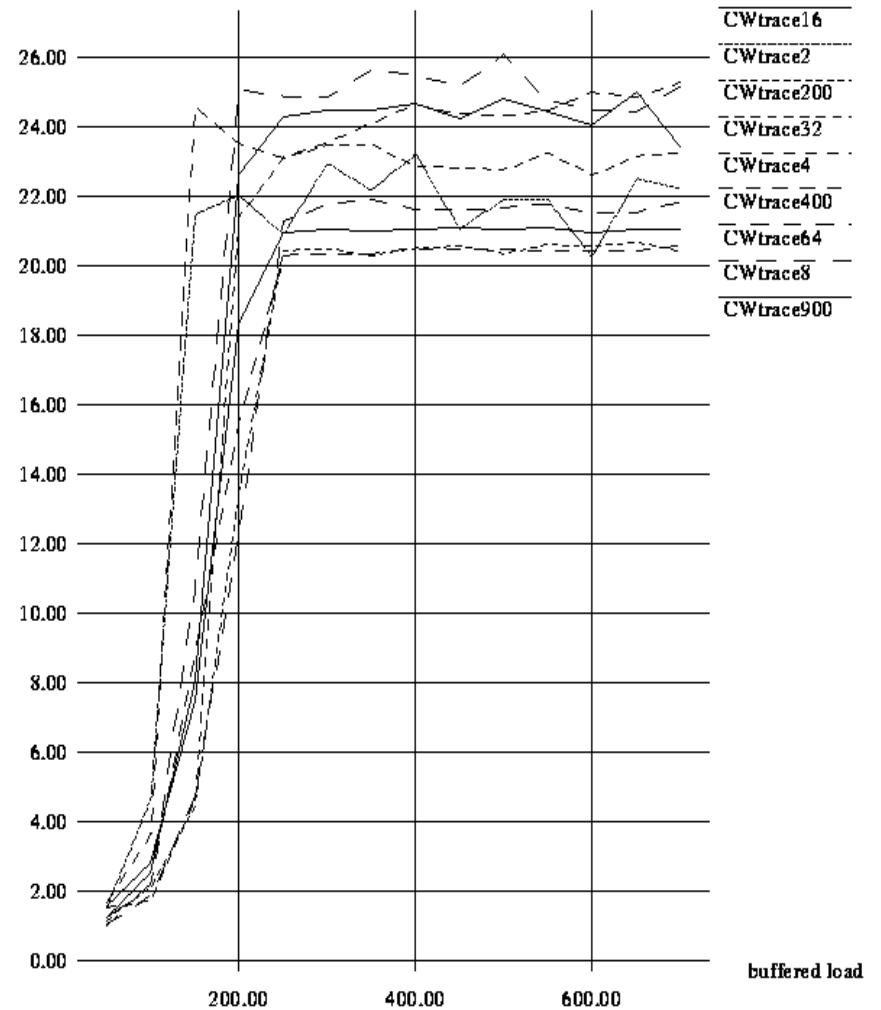
Throughput for different CW sizes, Leland trace

Throughput in kbyte/s x 10<sup>3</sup>



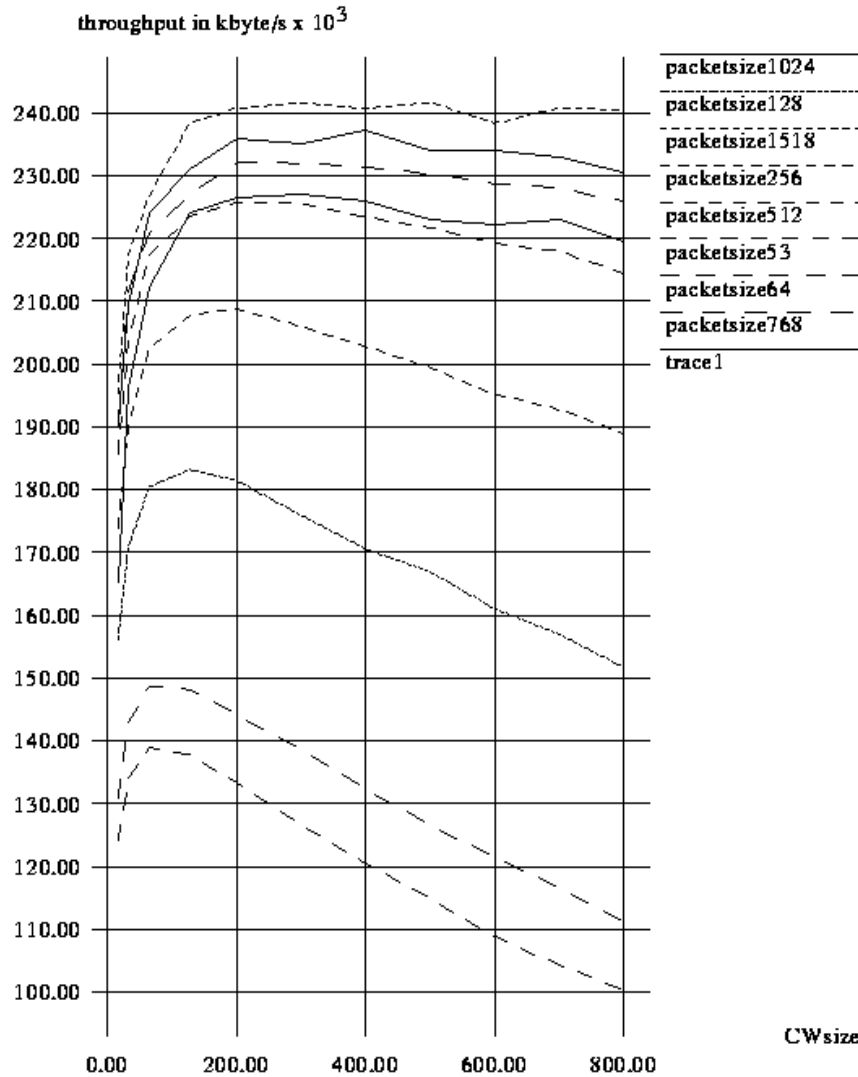
Access delay for different CW sizes, Leland trace

access delay in s x 10<sup>-3</sup>

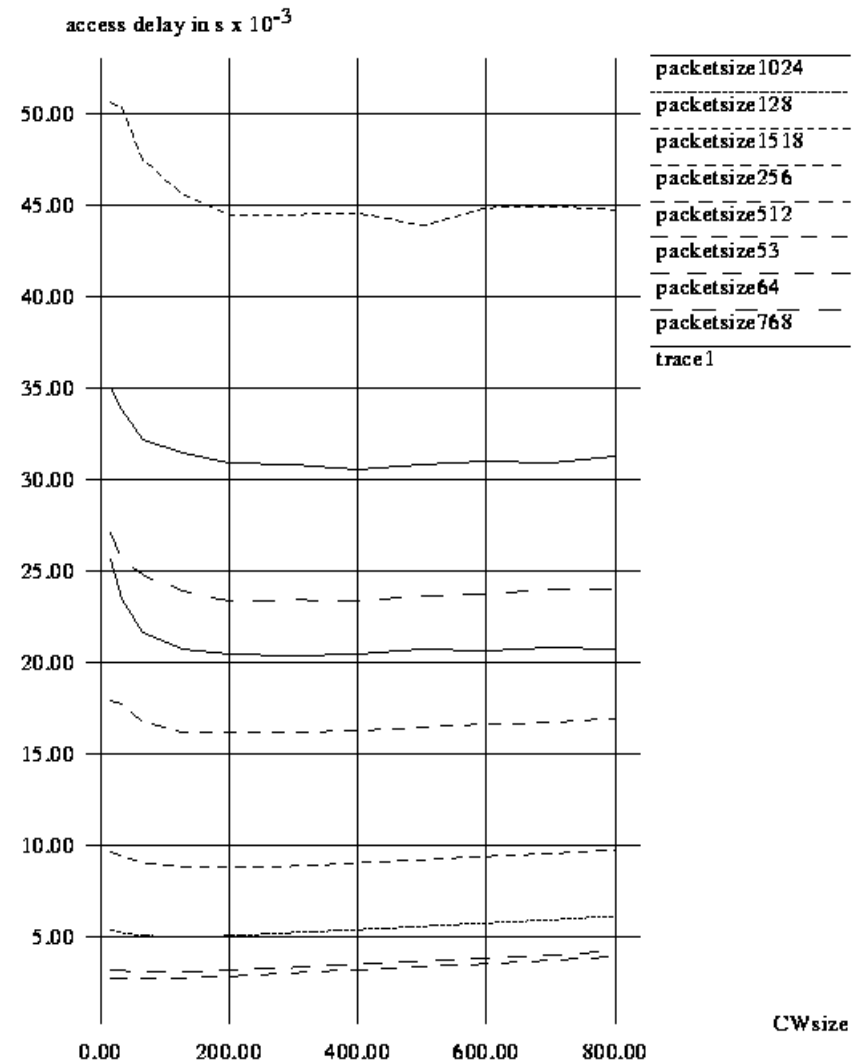


## Simulation results (cont.)

**Throughput for different packet sizes**

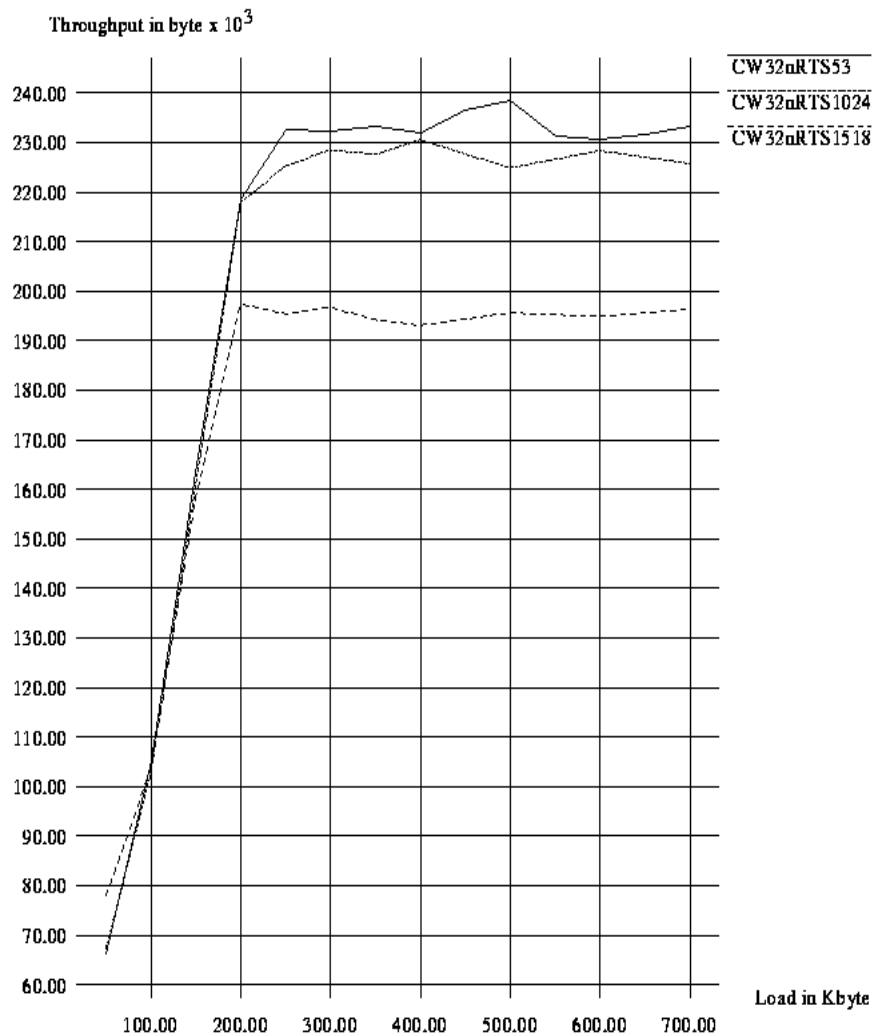


**Access delay for different packet sizes**

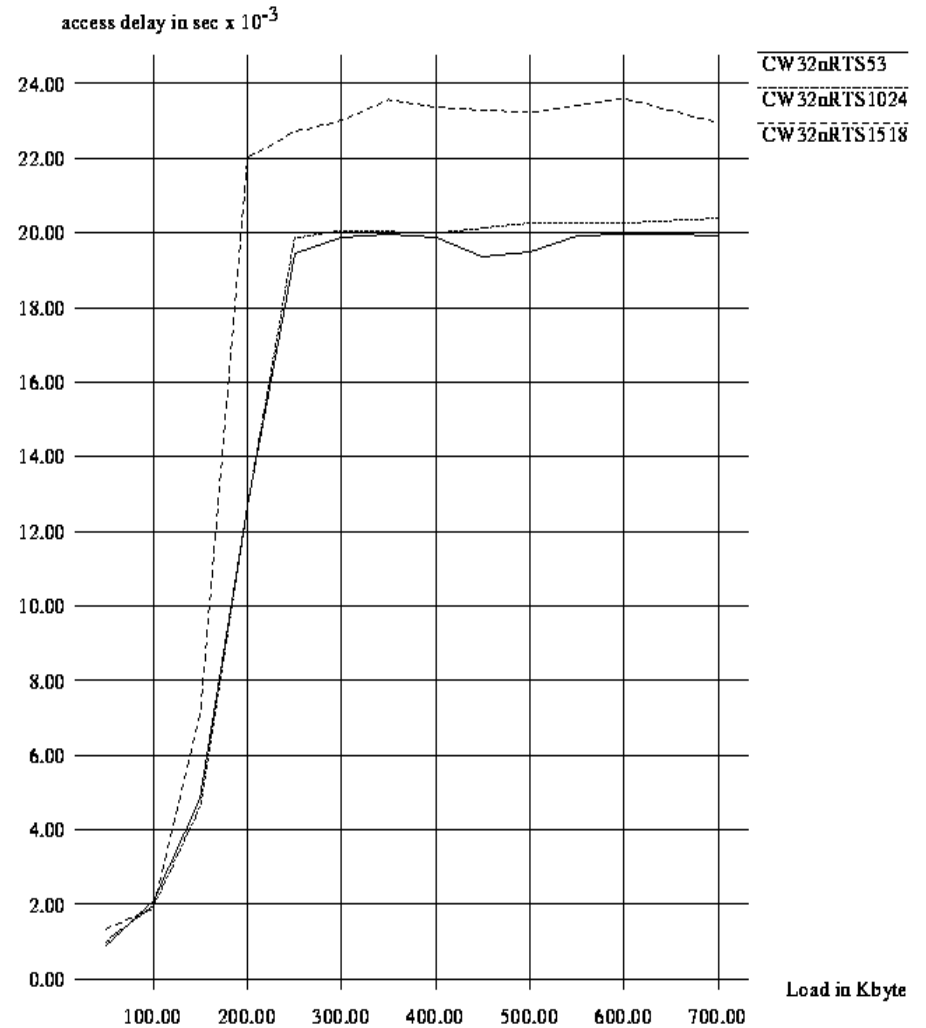


## Simulation results (cont.)

**DFWMAC with different NoRTS values, Leland trace, CW=32**

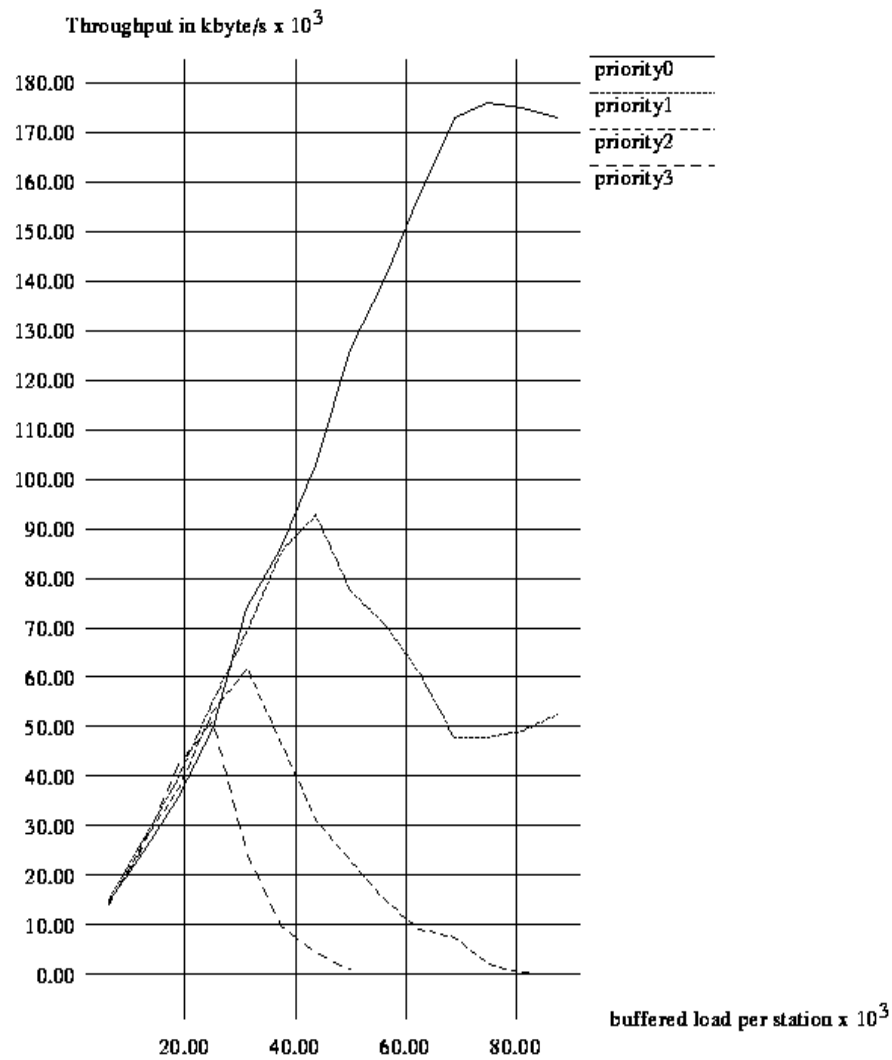


**DFWMAC with different NoRTS values, Access delay**

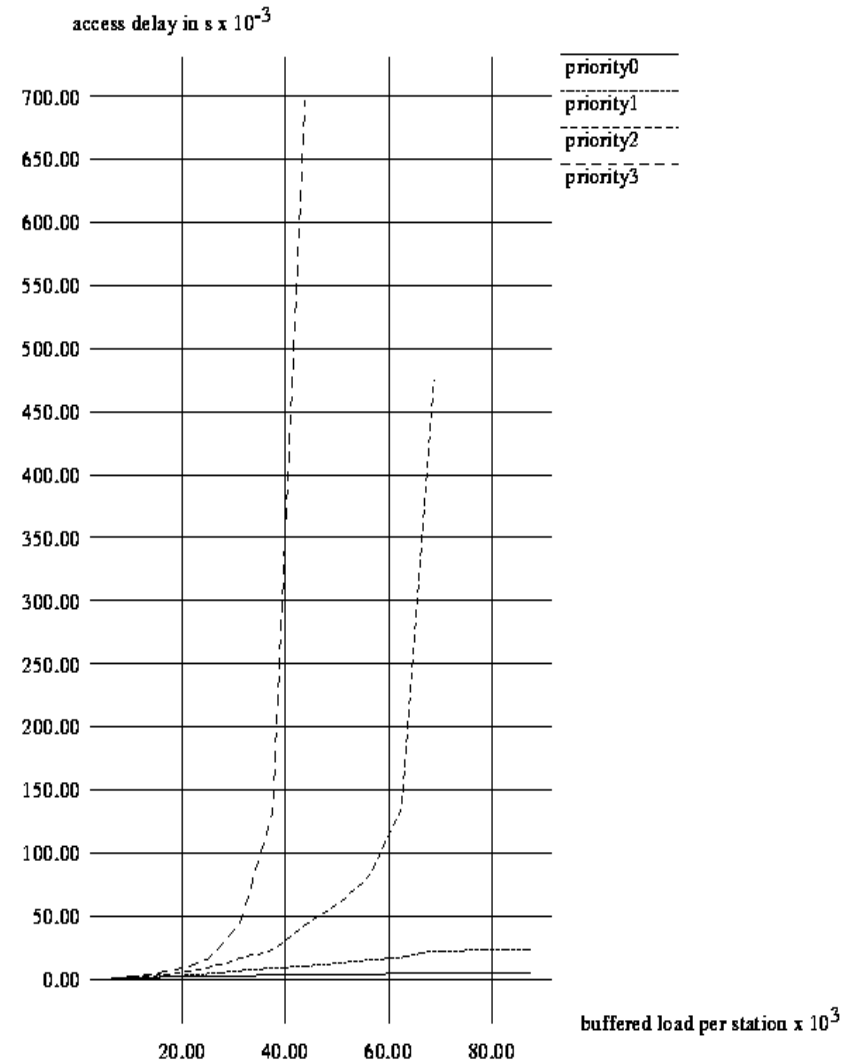


## Simulation results (cont.)

**Throughput for 4 priority classes (2 stations per class)**

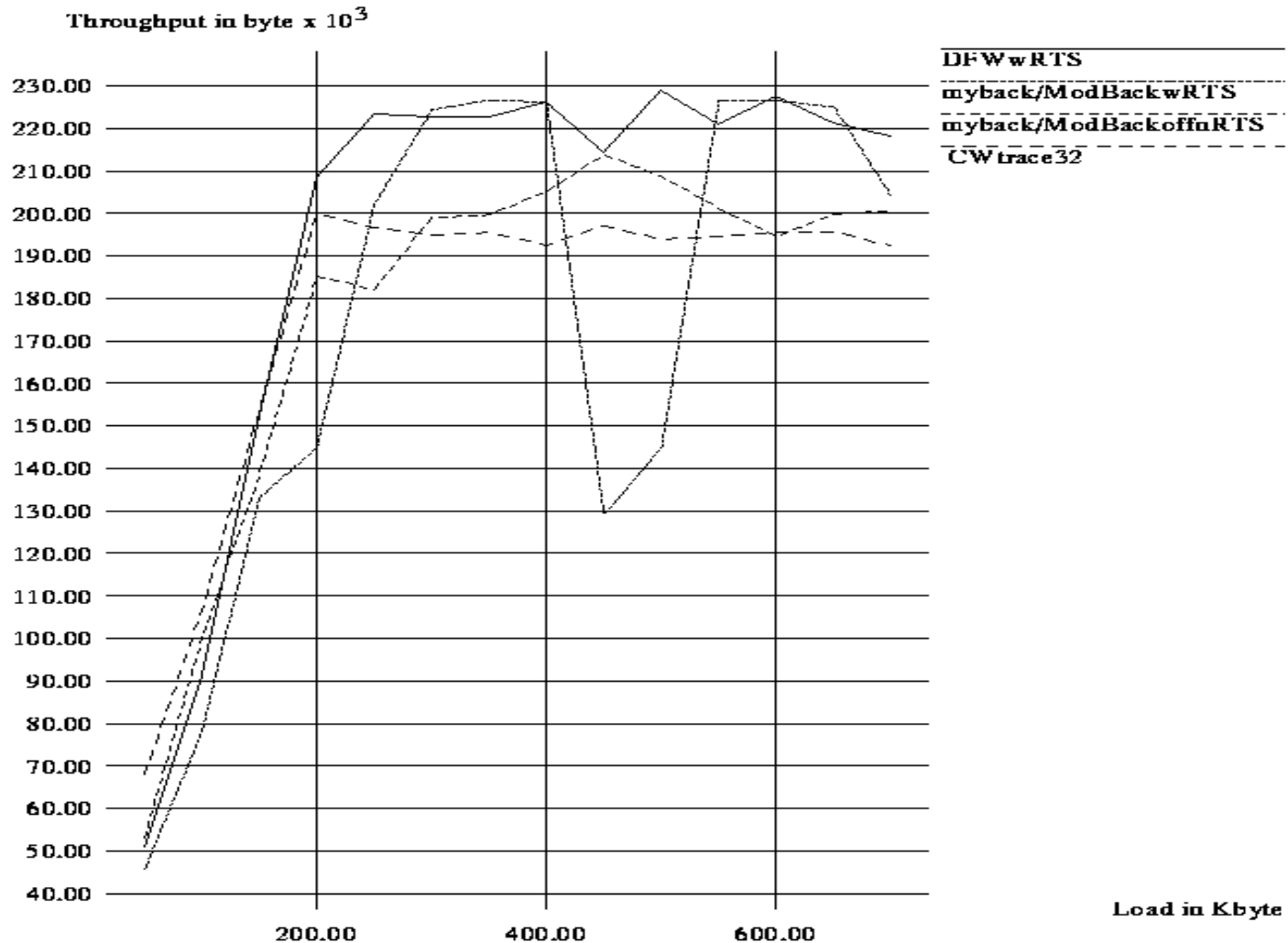


**Access delay for 4 priority classes (2 stations per class)**

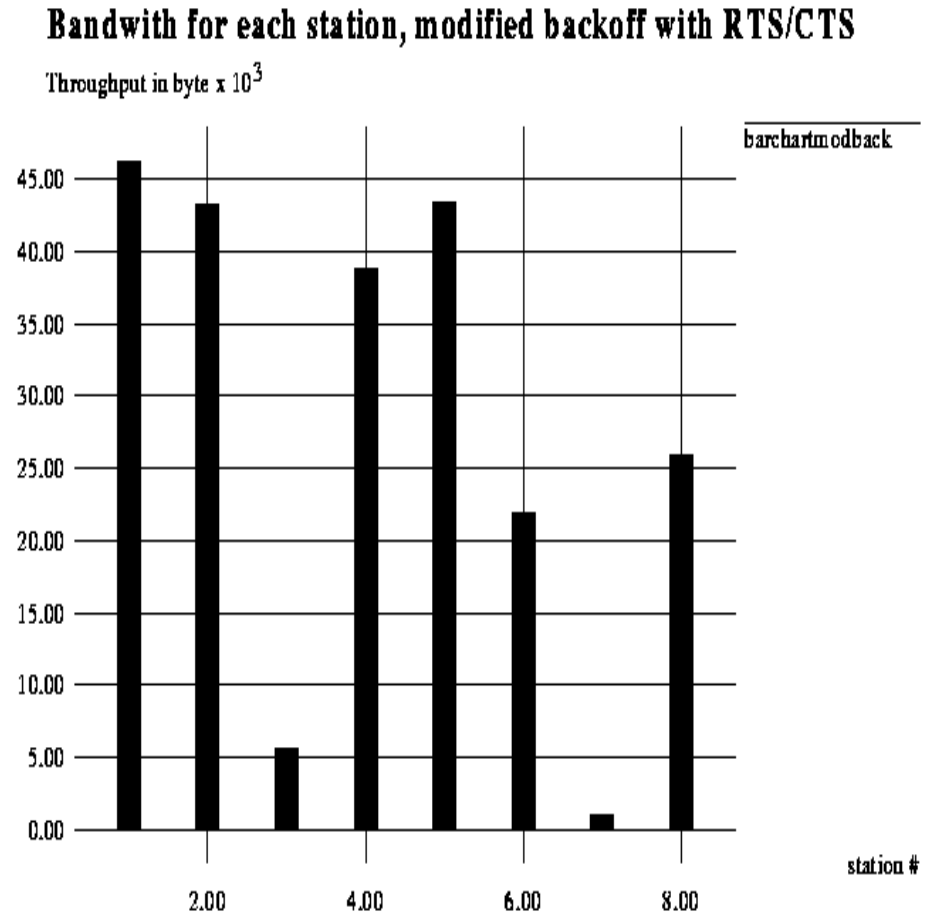
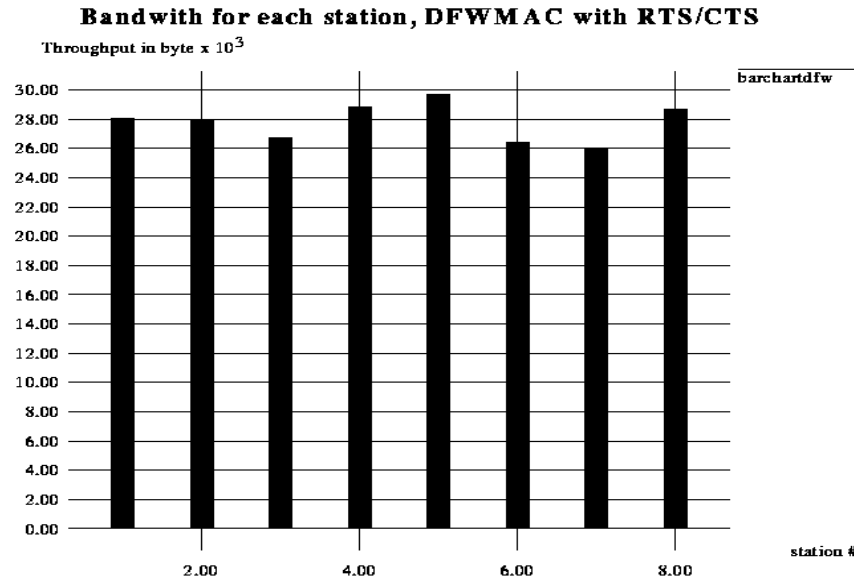


## Simulation results (cont.)

### Different protocol strategies, Leland trace, CW32



## Simulation results (cont.)



## *Conclusions*

- **Contention window size should depend on:**
  - ♦ *number of stations*
  - ♦ *packet size*
- **RTS/CTS leads to a better protocol behaviour even in the case of exposed stations**
- **Current backoff scheme has reserves**
  - ⇒ *modified backoff scheme leads to a better throughput but to unfairness*
  - ⇒ *more research on backoff schemes*

## **Near Future**

- **more detailed results on DFWMAC**
- **other protocols like URN, COMB, EY-NPMA, WATM**

## **Farther Future**

- **contributions to a High-speed WMAC (HWMAC)**